

Not Discrete Enough: On the Inherent Insecurity of dTPMs for Measured Boot

Cyber Security Experimentation and Test (CSET'25)

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TPM

- **Trusted Platform Module (TPM):** Hardware root of trust for cryptographic operations
 - Discrete TPM (dTPM): Separate chip, external communication bus
 - Firmware TPM (fTPM): Software implementation within CPU package

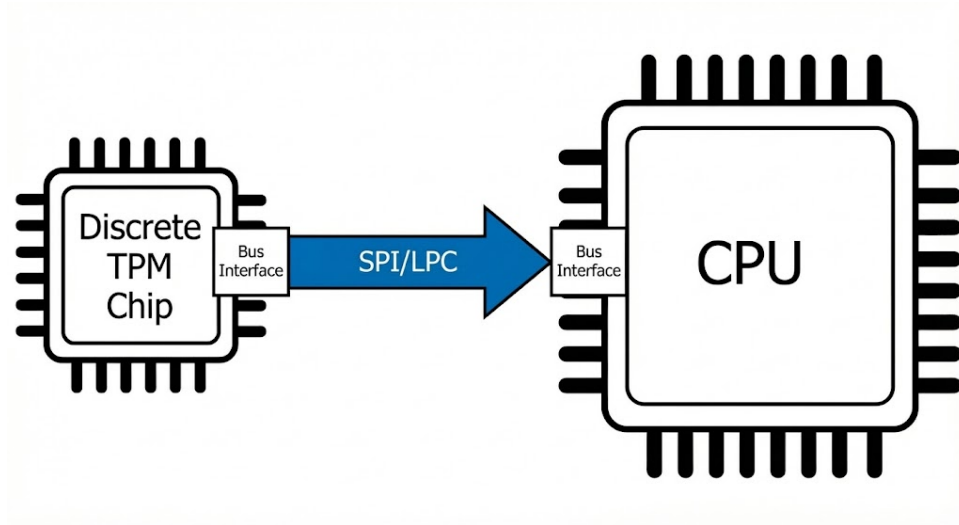


Fig: dTPM

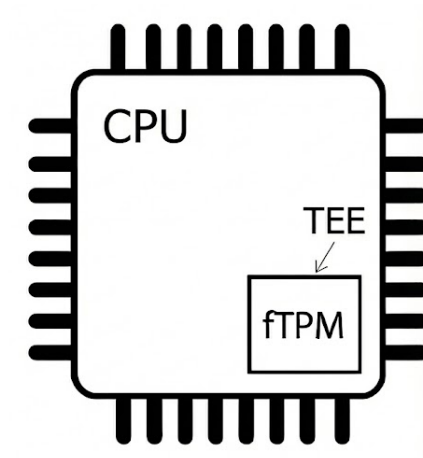


Fig: fTPM

Platform Configuration Registers & Full Disk Encryption

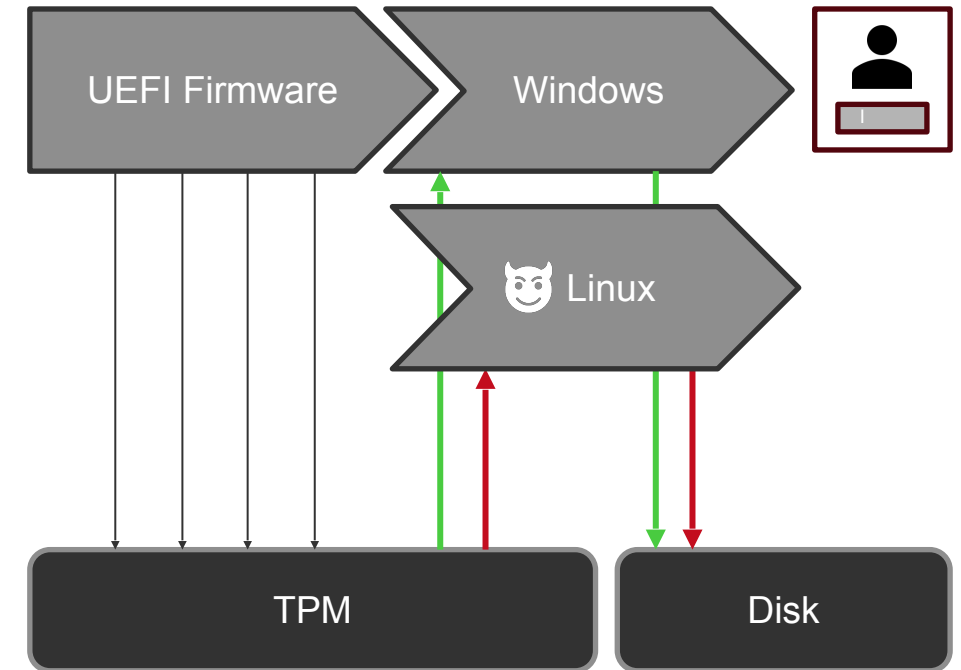
- **PCR:** Registers to store immutable states
- **Hash Chain:** Measuring system state with PCR_EXTEND
 - $\text{PCR_new} = \text{Hash}(\text{PCR_old} || \text{measurement})$
 - Reset to Zero only with TPM Reset
- **Full Disk Encryption:** BitLocker/LUKS seal keys to specific PCR states
 - PCR 7: Secure Boot state and authorities
 - PCR 11: Boot Manager "lock" mechanism (initially zeros, extended after unsealing)
 - TPM only protector

Attacks against discrete TPMs

- **Typical attack path:** Sniff bus communication of dTPMs to steal key
 - Could be mitigated through bus encryption techniques
- **Another security assumption:** TPM reset only occurs with platform reset
- **OSLO:** Bernhard Kauer performed a TPM (v1.2) reset attack 18 years ago ([16th USENIX](#))
- **Validation:** 2 dTPM platforms (Asrock B850 & Lenovo ThinkPad T480 with Infineon SLB9670)
- **Research Questions:**
 - Can the dTPMs (v2.0) on modern hardware still be reset from an attacker-controlled environment?
 - Can event logs be extracted and replayed to reconstruct BitLocker's expected PCR states?

Threat Model

- Goal: Protect laptop when it's stolen or lost
- Windows **Lockscreen** protects *running* system from unauthorized access
- However, Windows disk could be read from an attacker-controlled Live Linux
- With **TPM-backed** Full Disk Encryption and a key protected by the TPM:
Live Linux will not be able to access the key (sealed by TPM) or disk (encrypted)



Attack Overview - End to End Process



 **Total Time: ~10 minutes** •  **Cost: \$0–5 in basic tools**

UEFI Secure Boot vs. Measured Boot

1. Secure Boot

- UEFI Firmware enforces integrity ("Verify then Execute")
- **Chain of Trust:**
 - *Root (PK) → KEK → Signature DB (db) / Deny list (dbx)*
- **Outcome:** Blocks unauthorized bootloaders in UEFI
- **Commonly used certificates:**
 - Windows [...] CA → signs Windows bootloaders
 - Microsoft [...] CA → signs 3rd party bootloaders (e.g. Linux shim)

2. Measured Boot

- Firmware extends measurements of code and configuration into **TPM PCRs** before transferring control
- PCR7 records Secure Boot state and configuration
- **Outcome:** Immutable* evidence of the boot state

PCR7 Event Log Analysis

- ① **EV_EFI_VARIABLE_DRIVER_CONFIG** **SecureBoot**
Secure Boot status
Inside: 1-byte state for **SecureBoot** variable (enabled=1, disabled=0), plus GUID/attributes.
- ② **EV_EFI_VARIABLE_DRIVER_CONFIG** **PK**
Platform Key
Inside: root of trust public key material (X.509 in EFI signature lists) for Secure Boot.
- ③ **EV_EFI_VARIABLE_DRIVER_CONFIG** **KEK**
Key Exchange Keys
Inside: keys authorized to update **db** / **dbx** (typically OEM and Microsoft keys).
- ④ **EV_EFI_VARIABLE_DRIVER_CONFIG** **db**
Signature Database
Inside: authorized certificates and/or hashes for EFI images.
- ⑤ **EV_EFI_VARIABLE_DRIVER_CONFIG** **dbx**
Forbidden Signature Database
Inside: revoked certificates and blocked image hashes.
- ⑥ **EV_SEPARATOR** **0x00000000**
Transition delimiter
Inside: fixed 4-byte zero value marking the boundary before authority measurements.
- ⑦ **EV_EFI_VARIABLE_AUTHORITY** **db**
Verification authority
Inside: the actual certificate chain from **db** used to validate code (issuer/subject, etc.).

Secure Boot on or off?

Secure Boot configuration

Who authorized the current boot loader?
(→ is it a Windows or a 3rd party boot loader?)

Event Digest Reconstruction Challenge

- **Why it can't be replayed:**

EV_EFI_VARIABLE_AUTHORITY event measures the Secure Boot certificate used to validate bootloader

⑦ **EV_EFI_VARIABLE_AUTHORITY** db
 Verification authority
 Inside: the actual certificate chain from db used to validate code (issuer/subject, etc.).

- **Required Components:**

- Signature Owner GUID: Globally unique based on authority Windows/Linux
- Certificate: Microsoft Windows Production CA 2011

- **Technical Solution:**

- Reverse-engineered firmware digest calculation

UEFI_VARIABLE_DATA

VariableName (GUID)	16 bytes	d719b2cb-3d3a-4596-...
UnicodeNameLength	8 bytes	0x0002
VariableDataLength	8 bytes	cert_size + 16
UnicodeName	4 bytes	"db" (UTF-16LE)
VariableData	variable	EFI_SIGNATURE_DATA
SignatureOwner: 77fa9abd-0359-4d32-... (16 bytes) SignatureData: X.509 Certificate (DER format)		

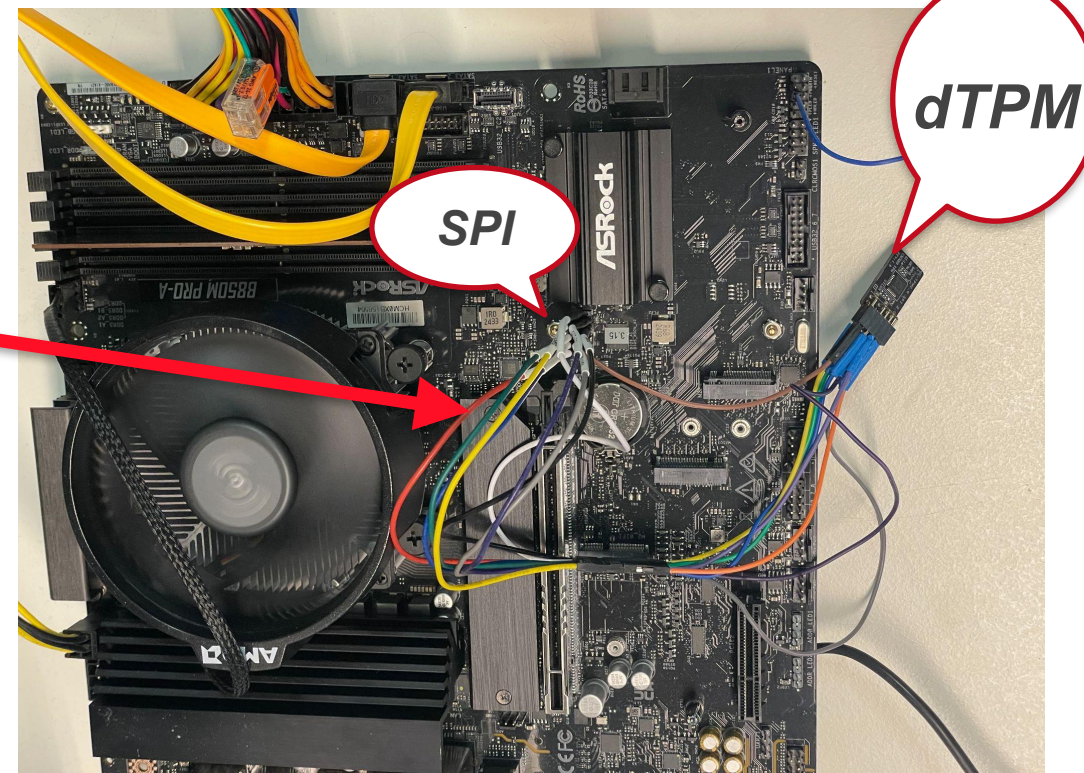
Complete PCR7 Replay Process

1. Parse event log into structured format
2. Replay PCR7 events up to EV_SEPARATOR (firmware/OS boundary)
3. Manual event digest calculation: Insert forged EV_EFI_VARIABLE_AUTHORITY Digest

Expected PCR 11 Digest is zero, so no replay required

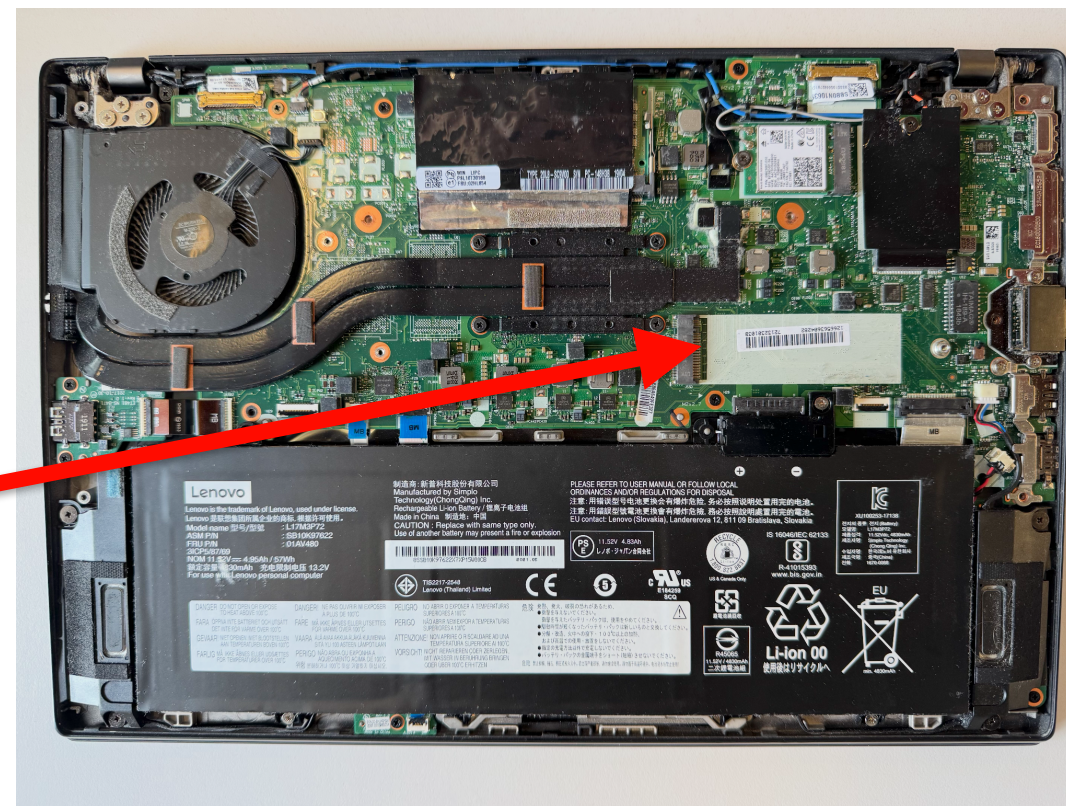
Hardware Reset - Asrock B850 + Infineon SLB9670

- **Challenge:** Reset TPM without host interruption
- **Platform Setup:** Dedicated dTPM module on development board
- **Reset Method:** VCC interruption Jumper Cable
- **Process:** Driver unbind → VCC interrupt → Driver rebind → Verify PCRs = 0
- **Result:** Clean Startup(CLEAR) achieved without host reboot



Hardware Reset - Lenovo ThinkPad T480 + Infineon SLB9670

- **Real-World Challenge:** Built-in dTPM, no exposed headers
- **Physical Access:** M.2 slot provides reset mechanism
 - Remove bottom cover
 - Access M.2 SSD slot
 - Ground PERST# pin on M.2 connector
- No full motherboard disassembly required



Hardware Reset - Lenovo ThinkPad T480

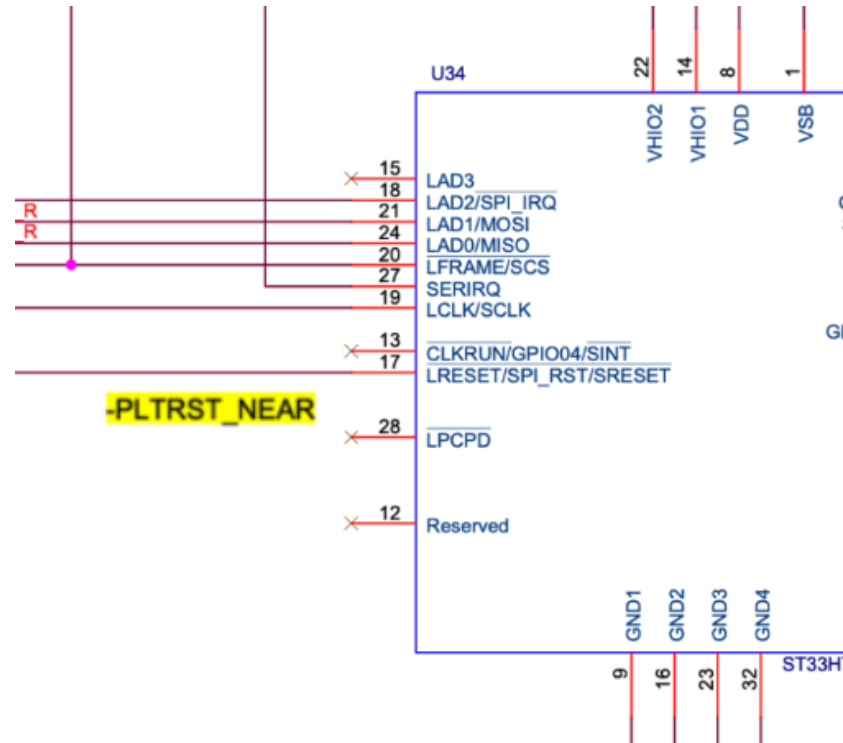
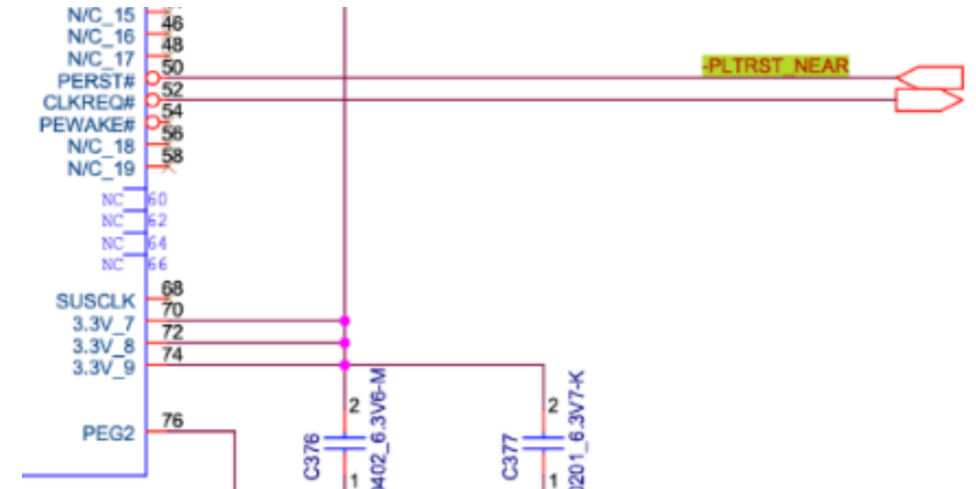


Fig: DTPM Schematic



M.2 Socket 3 (Key-M) for 2280 S3 SSD

Bitlocker VMK Unsealing

- **Sealed Object Extraction:**
 - Parse BitLocker metadata from volume header using dislocker
 - Locate Payload after line “`--> ENTRY TYPE VMK”
 - Extract TPM2B_PUBLIC and TPM2B_PRIVATE structures
- **Unsealing Process:**
 - Load sealed object
 - Start TPM policy session
 - Apply PolicyPCR for PCR 7,11 digest verification
 - Execute TPM2_Unseal command
 - Retrieve decrypted Volume Master Key

```

0x00000000 00 aa 00 20 c9 67 24 a2-4e 5c c9 83 55 5b 8d 56
0x00000010 0b fe af 3f ce 92 93 ff-a8 8c 7c 67 05 5e 74 01
0x00000020 eb 1d 48 eb 00 10 40 be-ee d3 1d 6e 36 8e 5c 08
0x00000030 1d a2 57 d3 1d f9 36 e0-4a 23 b4 10 b6 85 71 a4
0x00000040 f5 60 99 2c c9 0f 31 de-20 df 39 a8 7a de 29 3f
0x00000050 be 6e 61 5b b0 33 ec a0-09 55 39 3c 19 fa dd ad
0x00000060 e8 8f 6b 7b ca 80 61 5d-1f 95 92 c7 7a 0a af 74
0x00000070 af b8 78 e4 d0 f9 b0 0d-3f 42 c7 5c 58 d1 68 67
0x00000080 2b a2 62 10 a7 fe 29 6a-86 12 67 0e ea 9f 22 55
0x00000090 32 7a 10 03 00 99 44 ce-26 9d e0 72 81 21 d7 3e
0x000000a0 9d bc 91 34 d2 6c 28 50-73 13 6c e8 00 4e 00 08
0x000000b0 00 0b 00 00 04 12 00 20-71 ad 88 82 14 fb ed de
0x000000c0 ce c3 53 00 bd be f8 6b-28 71 d7 ee 38 89 58 54
0x000000d0 ed d5 5a d3 6e 77 7e fe-00 10 00 20 4d 28 54 cc
0x000000e0 3e 27 d5 6b 67 4b d0 be-f2 7e ba eb 00 b9 7f a3
0x000000f0 cf bb c2 a5 17 a7 7d 77-75 5a e6 9e 00 20 7f 67
0x00000100 36 18 e2 b6 f9 5d e8 4f-77 cc f8 ee bc d9 ca 3f
0x00000110 77 e2 a6 f2 f9 91 6c 6f-7e 29 e9 24 83 1a 03 80
0x00000120 08 00
  
```

TPM2B_PRIVATE

TPM2B_PUBLIC

Results

- **Validation:** Attack successfully works on both dTPM platforms
- Total attack time ~10 minutes
- **Costs:** Significantly cheaper than bus sniffing

Method	Cost	Complexity
Bus Sniffing	\$10-1,499	Medium
Reset/Replay	\$0-5	Easy

Conclusion

- **Countermeasure Effectiveness:**

- ✓ Strong: Firmware/Integrated TPM (eliminates external reset capability)
- ✓ Strong: TPM+PIN (adds unreplayable human factor)
- △ Limited: BIOS passwords
- ✗ Ineffective: Bus encryption (protects parameters but not event log)
- △ Effective but breaks Spec: PCR hiding. PCR transparency required for attestation

- **Recommendations:**

- TPM+PIN must be minimum standard for dTPM systems
- fTPM or other integrated TPM for new designs

Thank You

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